

## Question 1 [40%]

- (i) Write a fragment of Java using Map / ArrayBasedMap that declares and creates a map object named nums, populates the map with the numbers from 1 to 100 inclusive (as keys) and then removes all those numbers divisible by 17 (i.e. 17, 34, 51, ...). (8%)

### Answer (Not Answer but the Full MapArray Implementation)

```
import helpers.Comparator;
import helpers.IntegerComparator;
import helpers.MapEntry;

public class MapArray<KeyType, ValueType> implements Map<KeyType, ValueType> {

    protected static final int INIT_CAPACITY = 100;           /* the initial capacity of the map
    */
    protected int capacity;                                   /* the current capacity
of the map */
    protected int numEntries;
    protected MapEntry<KeyType, ValueType>[] entries;
    protected Comparator<KeyType> comparator;                /* comparator that defines equality between
keys */
    protected static final int NO_SUCH_KEY = -1;             /* value denoting unsuccessful search */

    public MapArray(Comparator<KeyType> myComparator) {
        entries = new MapEntry[INIT_CAPACITY];
        capacity = INIT_CAPACITY;
        comparator = myComparator;
        numEntries = 0;
    }

    public int size() {
        return numEntries;
    }

    public boolean isEmpty() {
        return (numEntries == 0);
    }

    public ValueType get(KeyType k) {
        int indexWithKey = findEntry(k);

        if (indexWithKey != NO_SUCH_KEY) {
            return (entries[indexWithKey]).getValue();
        }
        return null;
    }

    public ValueType put(KeyType k, ValueType e) {
        MapEntry<KeyType, ValueType> newEntry = new MapEntry<KeyType, ValueType>(k, e);
        int indexWithKey = findEntry(k);

        if (indexWithKey != NO_SUCH_KEY) {
            ValueType oldVal = entries[indexWithKey].getValue();
            entries[indexWithKey] = newEntry;
            return oldVal;
        }
        else {
            expandIfNecessary();
            entries[numEntries++] = newEntry;
        }
        return null;
    }
}
```

```

public ValueType remove(KeyType k) {
    int indexWithKey = findEntry(k);

    if (indexWithKey != NO_SUCH_KEY) {
        MapEntry<KeyType, ValueType> removedEntry =
            entries[indexWithKey];
        entries[indexWithKey] = entries[numEntries-1];
        numEntries--;
        return removedEntry.getValue();
    }
    return null;
}

/*****
 * Helper Methods
 *****/
private int findEntry(KeyType key) {
    for (int i = 0; i < numEntries; i++) {
        if ( comparator.compare(key, entries[i].getKey()) == 0 ) {
            return i;
        }
    }
    return NO_SUCH_KEY;
}

protected void expandIfNecessary() {
    if (numEntries == capacity) {
        int newCapacity = 2*capacity;
        MapEntry<KeyType, ValueType>[] temp = new MapEntry[newCapacity];

        for (int i = 0; i < capacity; i++) {
            temp[i] = entries[i];
        }
        entries = temp;
        capacity = newCapacity;
    }
}
}

```

---

### REMOVE NUMBERS DIVISIBLE BY 17 ???

```

public remove() {
    capacity = 100;
    int number = 1;

    while (number <= capacity) {
        if (number % 17 == 0) {
            nums.remove(number);
        }
    }
}

```

- (ii) Give a pseudocode fragment that takes the contents of a queue (ADT Queue) and reverses the order of the items contained within it. You may make use of an additional

ADT from {Stack, Queue, List, Map}, if you wish.

(11%)

Answer

Algorithm ReverseQueue(Q)

Input: queue Q

Output: queue Q in reverse order

S is an empty stack

```
while(!Q.isEmpty()) do
    S.push(Q.dequeue())
```

- (iii) The incomplete non-recursive binary search shown is intended to return the index within array S that contains the search key k (or “index” -1 if the search key is not present). Complete the algorithm by providing appropriate pseudocode for the placeholders labelled [a] etc. (8%)

Answer

Algorithm BinarySearch(S, k)

low <- 0

high <- S.size() - 1

while low < high do

mid = (low + high)/2

midKey = key(mid)

if k = midKey then

return mid

else

if k < midKey then

high <- mid - 1

else

low mid+1

return - 1

- (iv) Consider an implementation of ADT List using a left-justified array representation based on the following instance variable declarations, where EltType stands for the element type of the items in the list.

```
private EltType elements[];
private int numElts;
```

Give a complete implementation for the operation remove.

(8%)

Answer

```
public EltType remove(int numElts) {
    checkIndex(numElts);
    EltType retElt = entries[numElts];

    entries[numElts] = entries[numEntries-1];

    numEntries--;
    return retElt;
}
```

- (v) State the number of comparisons completed during the execution of the following algorithm when applied to an array X of length n. Justify your reasoning carefully. (8%)

```
Algorithm BS(X, n):
s ← 1
while s < n do
    curr = 0
    while curr < n-1 do
        next = curr + 1
        currElt = X[curr]
        nextElt = X[next]
        if currElt > nextElt then
            X[curr] = nextElt
            X[next] = currElt
            curr ← curr + 1
    s ← s + 1
```

Answer

## Question 2 [30%]

Give a complete Java implementation from scratch for an enhanced version of the traditional queue ADT that includes all the usual ADT operations as well as the following:

`delete(val)`: Remove from the queue all elements that equal the specified value; return the number of deletions made. Input: `EltType`; Output: `int`

Your implementation must respect the following conditions:

1. It must be based on the concept of a doubly-linked list and
2. It must be capable of housing elements of any (comparable) data type.

You do not need to provide code for an interface nor for a node class.

### Answer

```
import helpers.LLNode;

public class QueueDLinkedList<EltType> implements Queue<EltType> {

    // Declarations
    private int size;
    private LLNode<EltType> head;
    private LLNode<EltType> tail;

    public QueueDLinkedList() {
        size = 0;
        head = new LLNode<EltType>(null, null, null);
        tail = new LLNode<EltType>(head, null, null);
        head.setNext(tail);
    }

    public int size() {
        return size;
    }

    public boolean isEmpty() {
        return (size == 0);
    }

    public static void main(String args[]) {

        QueueDLinkedList<Integer> map = new QueueDLinkedList<Integer>();

        map.enqueue(1);
        map.enqueue(2);
        map.enqueue(3);
        map.enqueue(4);
        System.out.println("Front: " + map.front());
        System.out.println(map.dequeue());

        map.showMap();
    }

    public EltType front() {
        if (size == 0) {
            flagError("illegal queue op");
        }

        LLNode<EltType> firstNode = head.getNext();
        return firstNode.getElement();
    }
}
```

```

public void enqueue(EltType obj) {
    LLNode<EltType> oldLast = tail.getPrev();
    LLNode<EltType> newLast = new LLNode<EltType>(oldLast, tail, obj);

    oldLast.setNext(newLast);
    tail.setPrev(newLast);

    size++;
}

public EltType dequeue() {

    if (size == 0) {
        flagError("illegal queue op");
    }

    LLNode<EltType> oldFirst = head.getNext();
    //head = head.getNext();
    head.setNext(oldFirst.getNext());

    size--;
    return oldFirst.getElement();
}

/*****
 * Helper Methods
 *****/
public void showMap() {
    System.out.println("\n****Start Map Structure****");
    System.out.println("Headnext: " + head.getNext().getElement());

    // get starting node
    LLNode<EltType> currentElement = head.getNext();

    // loop through map
    for(int i = 0; i < size; i++) {

        System.out.println("element at position " + i + ": value = " +
currentElement.getElement());

        // go to next node
        currentElement = currentElement.getNext();

    }

    System.out.println("Tailprev: " + tail.getPrev().getElement());
    System.out.println("****End Map Structure****");
}

private void flagError(String errmsg) {
    System.out.println("LinkedQueue: "+errmsg);
    System.exit(1);
}
}

```

### Question 3 [30%]

- (i) The following recursive algorithm segregates the contents of A[f...r] (the segment of array A between indices f and r inclusive) so that all the values less than or equal to x

appear to the left of those greater than x.

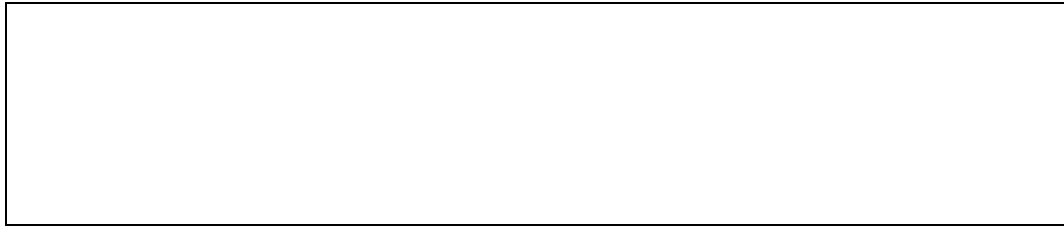
```
Algorithm Split(A, x, f, r):
  if r < f then
    return
  else
    if A[f] ≤ x then
      Split(A, x, f+1, r)
    else
      if A[r] > x then
        Split(A, x, f, r-1)
      else
        temp = A[f]
        A[f] = A[r]
        A[r] = temp
        Split(A, x, f+1, r-1)
```

Draw a recursion tree to show the execution of Split (X, 5, 0, 7) where the array contains [3, 8, 4, 1, 6, 5, 2, 7] initially. Show the state of the array and the values of f and r at each stage. (6%)

Answer

- (ii) Argue that for any array A of length n, the number of calls to Split arising from Split(A, x, 0, n-1) for any x is at most n + 1 (6%)

Answer



- (iii) Show how the algorithm may be modified to produce a variant Split2 so that Split2(A, x, f, r) not only partitions the elements in the manner of Split but also returns the number of values in A[f...r] that are less than or equal to x. (9%)

Answer (Worth 3 / 9 as is)

```
Algorithm Split2(A,x, f, r)
  count ← 0
  if r < f then
    return
  else
    if A[f] ≤ x then
      Split (A, x, f+1, r)
      count++
    else
      if A[r] > x then
        Split (A, x, f, r-1)
      else
        temp = A[f]
        A[f] = A[r]
        Split (A, x, f+1, r-1)
```

- (iv) Based on the Split2 algorithm, write a recursive sorting algorithm (in pseudocode) that takes an interval of an array A[f...r] and that re-arranges the contents so that they appear in increasing order left to right. (9%)

Answer





